

Workload Characterization of SPECweb2005

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As the diversity of computer applications increases, so does the need to come up with newer benchmarks. Moreover, as the usage characteristics of applications change, it becomes necessary to update the relevant benchmarks to reflect these characteristics in the benchmarks appropriately. SPECweb96 was the first web benchmark based on the characteristics of ISPs that downloaded only static pages. This was replaced with SPECweb99 in 1999. The main difference between SPECweb96 (see [1]) and SPECweb99 (see [2]) was the inclusion of simple dynamic pages in SPECweb99. While this was a major step forward, the standard web benchmark still suffered from several shortcomings (see [3] for a discussion and analysis), from a customer benchmarking point of view, the main one being that SPECweb99 did not include SSL based requests. With this in mind, SPECweb99_SSL was created by adding an SSL envelope around SPECweb99 (for details, see [4]). Together, the two benchmarks served the community for

more than three years. However, it was clear that the benchmarks did not represent any real world web workload. The effort towards SPECweb2005 was focussed on representing specific web workloads with a hope to replicate and represent the individual characteristics from real world applications in this benchmark. Unlike its predecessors which were based on request models for files, SPECweb2005 is based on downloading pages, which are dynamically created files, followed by requests for embedded images within the files, with the page download considered as being complete when the dynamic page as well as all the embedded files are downloaded. Currently, this benchmark includes three different workloads:

- Banking, which is fully SSL based workload
- Ecommerce workload, which includes both SSL as well as non-SSL requests.
- Support workload, which is based fully on http requests and includes the download of really large files.

The goal of this paper is to summarize the characteristics of this benchmark, critique it to a limited extent and to present a basis for further commentaries and comparisons

which could pave way to the improvement of the benchmark as technology changes and new usage patterns emerge.

Characteristics of Banking workload

This workload is based on internet personal banking. Each individual logs into the system and checks for their account status, followed by either bill payment, transfer of money between accounts, or changing their personal profile. Operations include both Post and GET operations. The system maintains a user session id for each user that logs in. About 20% of the incoming users do not logout, letting the session time out. Each request session consists of processing a dynamic request followed by requests for several static images. The static image requests are a function of the page requested. Some of the static requests are returned with a “304” response. The dynamic request is first sent over a TCP connection, then another TCP connection is initiated for requesting the static files. The static files are thus requested over the two TCP connections. The second TCP connection uses the same SSL session as the first one. The SSL session is reused for all subsequent requests within the session, until the user logs off or abandons

the session. After each page request, the user thread goes through a Think Time, which averages about 9.98 sec.

We will now present some of the workload characteristics that are derived from the workload design.

We refer to the main metric for this workload as SPECweb2005_Banking. This is the number of simultaneous active user threads supported while running the workload. We will refer to each page request operation as SW2005_Op. So, each SW2005_OP will include one or more http ops involving dynamic processing followed by a number of http static file requests.

- Average number of SW2005_Ops in a login session = 4.6 (derived from the Markov chain for the workload).
- Number of TCP connections made per login session = 2
- Number of Page requests per TCP connection = 2.3
- Average number of HTTP requests per SW2005_Op = 12 (average derived from design parameters)
- So, the average number of HTTP requests per TCP connection = $12 * 2.3$

= 27.6.

Using the Think Time value, average response interval seen from the system at peak load conditions, and the probability for the only state that does generate three page requests (for details, please see [5]), we can compute the average number of SW2005_Ops/s for each user session supported to be approximately 0.149. Thus, for every 1000 user sessions supported (SPECweb2005_Banking value of 1000), we get 149 SW2005_Ops/sec.

When the first user logs in, there is a full SSL handshake over the first TCP connection opened. The second TCP connection opened reuses the SSL session from the first one. In steady state, 21.47% of all requests are logins. Therefore, we get 32 full handshakes (149*0.2147) for every 1000 user sessions supported.

There are 2 TCP handshakes (accepts) for every user session. Given that there are an average of about 4.6 page requests per user session, 43.47% (2/4.6) of all SW2005_Ops will initiate a new TCP connection. As explained above, we have 149 SW2005_Ops/sec for every 1000 SPECWeb2005_Banking user sessions

supported, so this will result in a total of 65 TCP accepts/sec for every 1000 SPECweb2005_Banking.

Each SW2005_OP results in an average page size of 30.4 KB. Given 4.6 SW2005_OPs per user session, over an average interval of 30.9 sec (computed using the average Think Time, and average response interval), we get 0.1488 SW2005_Ops/sec per SPECweb2005_Banking session supported. This is equivalent to 4.52 KB/sec of Outgoing bytes per SPECweb2005_Banking session. For a score of 1000 SPECweb2005_Banking, this translates to 36 Mb/s of outgoing traffic from SUT to client.

The SUT also sends requests to the BackEnd Simulator (BeSim) and receives responses back. The following diagram shows the relative sizes of byte traffic between the client, SUT and the BeSim. These values are based on the measurements collected on our testbeds.

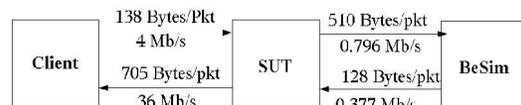


Figure 1 Network IO for Banking workload

<i>Summary characteristic</i>	<i>Value</i>
SW2005_Op/sec per 1000 SPECweb2005_Banking user sessions	149
HTTP- ops/s per 1000 SPECweb2005_Banking user sessions	1788 (=149*12)
SSL full handshakes/s per 1000 SPECweb2005_Banking user sessions	32
SSL resumed handshakes/sec per 1000 SPECweb2005_Banking user sessions.	32
TCP accepts/sec per 1000 SPECweb2005_Banking user sessions	65
Client to SUT traffic/sec per 1000 SPECweb2005_Banking user sessions	4 Mb/s
SUT to client traffic/sec per 1000 SPECweb2005_Banking user sessions	36 Mb/s
SUT to BeSim traffic/sec per 1000 SPECweb2005_Banking user sessions	0.796 Mb/s
BeSim to SUT traffic/sec per 1000 SPECweb2005_Banking user sessions	0.377 Mb/s
Average bytes/Pkt (Client to SUT)	138 Bytes

<i>Summary characteristic</i>	<i>Value</i>
Average bytes/pkt SUT to client	705 Bytes
Average bytes/pkt SUT to BeSim	517 Bytes
Average bytes/pkt BeSim to SUT	128 Bytes

Table 1: Banking workload characteristics

Other Characteristics and Recommendation

The only images that scale with the workload are the customer check images.

The workload scaling is mostly limited by CPU power. Memory and disk usage is not as heavy for this workload. Given that the resource usage characteristics are very platform dependent, we do not present these characteristics in this paper.

As the workload is scaled, BeSim is likely to present a bottleneck. The next release of the benchmark is expected to address this issue.

Characteristics of the Ecommerce Workload

The Ecommerce workload is based on the workload characteristics of a site that sells computers and related accessories. Part of this workload uses plain http based requests, while the other part uses https

based requests. A customer visiting the site passes through three distinct phases. The first phase is when the customer browses the website, looking for a product. This phase involves searches and browsing activity over product web pages, all of which are dynamically created. The second phase is the customization phase, where the user customizes the product. Finally, if the customer wants to check-out/buy the product, then she/he does that using SSL based requests. The workload uses 11 distinct scripts. As in the case of Banking workload, the transition between states is driven via a Markov Chain. Each page request (dynamic) is accompanied by a bunch of static requests for embedded images, including page images and product images. The product images are scaled with the benchmark.

From the benchmark design, we can compute the average number of SW2005_Ops for each incoming user session to be 8.8 and the number of http-ops per SW2005_OP as 17. As in the case of Banking workload, the http requests for each session uses 2 TCP connections. Using an average of 9.98 sec + 2 sec (for average response time under heavy load, since 3 sec is the QOS limit for

Time_Good [5][6]) between subsequent SW2005_OP, we can estimate the SW2005_Ops/sec for each user session supported as 0.093. Thus, for every 1000 sessions supported, we can expect to see 93.7 SW2005_OPs/sec and a corresponding 1580 HTTP ops/s.

Unlike in the case of the Banking workload, only a fraction of the user sessions enter the SSL stage. The fraction of the sessions that enter the checkout stage is about 2/3 (computed from the design Markov Chain). For each session, from start until exit, there will be exactly one full SSL handshake and one partial SSL handshake. Each user session involves 2 TCP connections if it does not enter the check out stage and 4 TCP connections if it enters the checkout stage. Using the above along with the information for average think-time and number of SW2005_Ops/user session, we can estimate the average the number of TCP connections/sec for every 1000 SPECweb2005_Ecom supported as 35.

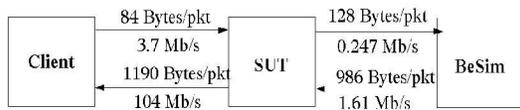


Figure 2 Network IO for Ecommerce

The following table summarizes these characteristics:

Characteristic	Value
SW2005_OPs/sec per 1000 SPECweb2005_Ecom	93
HTTP-OP/sec per 1000 SPECweb2005_Ecom	1580
SSL full handshakes/sec per 1000 SPECweb2005_Ecom	17
SSL resumed handshakes/sec per 1000 SPECweb2005_Ecom	17
TCP accepts/sec per 1000 SPECweb2005_Ecom	35
Client to SUT traffic/sec per 1000 SPECweb2005_Ecom	3.7 Mb/s
SUT to client traffic/sec per 1000 SPECweb2005_Ecom	104 Mb/s

Characteristic	Value
SUT to BeSim traffic/sec per 1000 SPECweb2005_Ecom	0.247 Mb/s
BeSim to SUT traffic/sec per 1000 SPECweb2005_Ecom	1.61 Mb/s
Average bytes/pkt (Client to SUT)	84 Bytes
Average bytes/pkt (SUT to Client)	1190 Bytes
Average bytes/pkt (SUT to BeSim)	128 Bytes
Average bytes/pkt (BeSim to SUT)	986 Bytes

Table 2: Ecommerce workload Characteristics

Other Characteristics and recommendations

In this workload, the product images scale linearly as the workload is scaled. Given that the IO needs for this workload are higher than those for Banking, memory and Disk IO play a significant role on the performance of this workload. Moreover, the bytes exchanged with the BeSim server and the number of BeSim queries for this workload are much higher than for the Banking workload. Given this, the BeSim scaling change is most important for this workload.

Characteristics of the Support Workload

The support workload was developed based on the characteristics seen in sites that were used to download upgrade and fix related patches for computer support. Typically, these downloads are very very large and the intent of this workload is to stress the IO and network IO. This workload does not use SSL. The user emulated in this workload, enters the system, searches for the right patch file to download, and then proceeds to download the file. The maximum size of the file downloaded is 35 MB. While the QOS for all other pages is based upon the total response interval seen, the QOS for the patch download page is based only on the throughput. 95% of the downloads are expected to meet a minimum bit-rate of 100 kb/s, and 99% are expected to meet a bit rate of 95 Kb/s. Also, the Think Time values used for this workload are about half as much as those for the other two workloads (5 seconds by design, 5.1 seconds in the implemented version).

From the design parameters, we can estimate the average number of page requests that an incoming user will make as 14.5. With an average Think Time

value of 5 sec between requests, and response time of about 1 sec for all requests except the one for the download file, and an average response time for the download file at about 68 seconds (measured), the cycle time comes to about 150 sec. Equivalently, dividing the cycle over 2 TCP connection requests, we get an average of 13 TCP connection requests/second, for each user emulated.

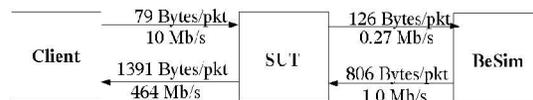


Figure 3 Network IO for Support workload

The following table summarizes some of the characteristics of this workload:

Characteristic	Value
SW2005_OPs/sec per 1000 SPECweb2005_Support	97 SW2005_Ops/sec
HTTP-Ops/sec per 1000 SPECweb2005_Support	2044 HTTP ops/sec
SSL full handshakes/sec per 1000 user sessions	N/A
SSL resumed handshakes/sec per 1000 user sessions	N/A

<i>Characteristic</i>	<i>Value</i>
TCP accepts/sec per 1000 user sessions	13
Client to SUT traffic/sec per 1000 user sessions	10 Mb/s
SUT to Client traffic/sec per 1000 user sessions	464 Mb/s
SUT to BeSim traffic/sec per 1000 user sessions	0.27 Mb/s
BeSim to SUT traffic/sec per 1000 user sessions	1.0 Mb/s
Average bytes/pkt (client to SUT)	79 bytes/pkt
Average bytes/pkt (SUT to client)	1341 bytes/pkt
Average bytes/pkt (SUT to BeSim)	126 bytes/pkt
Average bytes/pkt (BeSim to SUT)	806 bytes/pkt
Average Mbps for every 1000 connections supported (SUT-client)	464 Mbps

Table 3: Support workload Characteristics

Other Characteristics and Comments

This workload is highly IO intensive. The workload performance is significantly influenced by the Zipf parameter used. Currently the benchmark uses a Zipf parameter of 1.2. The effect of this

parameter is to determine the access pattern among the files. The higher the value, most of the access will be restricted to a minimal set of files. For the same file base, keeping the zipf parameter value higher would hence reduce the disk access on a relative basis.

We conducted a simple simulation to track the number of times a new file will be loaded into the memory, resulting in the purging of the existing files from the memory. We used a very simple memory model. First we assume that 100% of the memory will be dedicated toward serving the file system. Directory and file choices used in the model were as per specs for the Support workload in SPECweb2005. We start the memory with no file loaded in it. Each time a new file is read, the file is loaded into the memory. If the file was previously accessed and is already in the memory, the new file will be loaded from the memory. When the memory is fully loaded, and then encounters a file not loaded into the memory, the new file will be loaded while a file that has not been accessed for the longest time will be purged (if necessary, more than one file may be purged in order to load the new file). If this happens, it is called memory churn and we note the percentage of file

accesses that result in a memory churn. The following table shows the increase in memory churn values, as the zipf parameter is increased, for various directory sizes and memory sizes.

<i>Number of Director ies</i>	<i>Memory Size</i>	<i>Memory churn % for alpha = 1.0</i>	<i>Memory churn % for alpha = 1.2</i>
500	4 GB	4.95%	3.89%
	8 GB	4.19%	3.23%
	16 GB	2.94%	2.09%
	32 GB	0.00%	0.00%
1000	4 GB	8.62%	6.40%
	8 GB	8.19%	5.18%
	16 GB	6.80%	4.46%
	32 GB	4.23%	2.07%
2500	64 GB	0.00%	0.00%
	32 GB	10.76%	5.30%
5000	64 GB	6.14%	0.13%
	64 GB	10.98%	2.78%
	128 GB	1.90%	0.00%

Table 4: Memory Churn variation with zipf value used in Support workload

Further study on the effect of this parameter and a check on how it matches similar real world applications is certainly in order.

References

[1] An Explanation of the SPECweb96 Benchmark,

www.spec.org/osg/web96/webpaper.html

[2] SPECweb99 Release 1.02 whitepaper, www.spec.org/web99/docs/whitepaper.html

[3] E. M. Nahum, Deconstructing SPECweb99, Proceedings of 7th International Workshop on Web Content Caching and Distribution, August 2002.

[4] SPECweb99_SSL Design Document and Implementation Overview, www.spec.org/web99ssl/docs/whitepaper.html

[5] SPECweb2005 Design Document, www.spec.org/web2005/design.html

[6] SPECweb2005 Run Rules Document, www.spec.org/web2005/run_rules.html